
A Compositional Semantics for Modelica-style Variable-structure Systems

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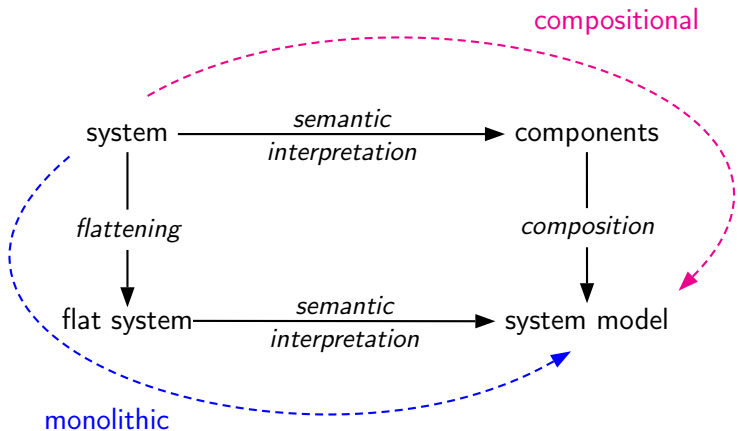


Goals

- Compositional semantics
 - readability and simplicity
 - separate compilation
 - variable-structure systems (structure dynamics)
- Separation of concerns
 - "ideal semantics"
 - "solver semantics"

By contrast, Modelica has a monolithic, flattening-based semantics with a mixture of conceptual and numerics-oriented aspects.

① *Compositional Semantics*



Ideal Semantics

Ideal Semantics

Ideal semantics ...

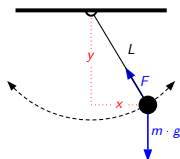
- lives in the realm of pure mathematics
- focuses on modeling concepts
- ignores numerical problems

Two levels of modeling concepts:

- **fixed-structure systems** (classical Modelica)
 ↪ simplify semantic presentation
- **variable-structure systems** (extended Modelica)
 ↪ prepare extended concepts (dynamic systems, separate compilation)

Ideal Semantics

Base Case: Atomic Fixed-Structure Component

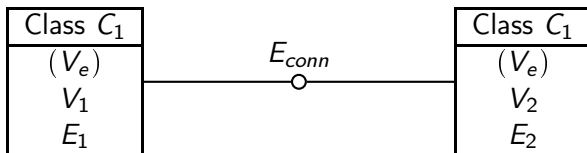


Class <i>Pendulum</i>	
PARAM <i>Length L</i>	parameters
PARAM <i>Mass m</i>	
CONST <i>Acceleration g</i>	
<i>Length x</i>	variables <i>V</i>
<i>Length y</i>	
<i>Force f</i>	
$m \cdot \ddot{x} = -\frac{x}{L} \cdot f$	equations <i>E</i>
$m \cdot \ddot{y} = -m \cdot g - \frac{y}{L} \cdot f$	
$x^2 + y^2 = L^2$	

Syntax:Class $C = (V_e, V_l, E)$ V_e : external variables V_l : local variables E : hybrid DAEsSemantics:Model $M = (F_e, F_l)$ F_e set of functions ($\hat{=} V_e$) F_l set of functions ($\hat{=} V_l$) $M \models E$

Ideal Semantics

Composition



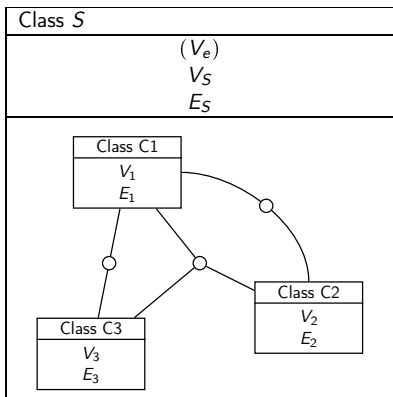
Composition is essentially a “pushout”:

$$\begin{aligned}
 \text{Mod}(S) &= (\text{Mod}(C_1) \otimes \text{Mod}(C_2)) \mid E_{conn} \\
 &\stackrel{\text{def}}{=} \{ M_1 \cup M_2 \mid M_1 \in \text{Mod}(C_1), \\
 &\quad M_2 \in \text{Mod}(C_2), \\
 &\quad M_1|_{V_e} = M_2|_{V_e}, \\
 &\quad M_1 \cup M_2 \models E_{conn} \}
 \end{aligned}$$

Note: Name clashes avoided by scoping rules (compiler)

Ideal Semantics

Subsystems

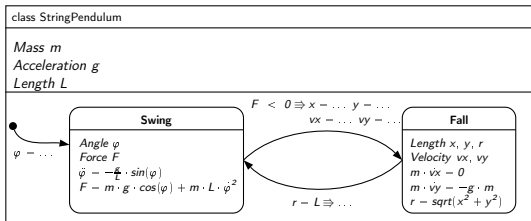
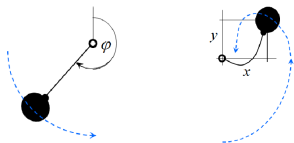


$$\text{Mod}(S) = (\text{Mod}(C_1) \otimes \cdots \otimes \text{Mod}(C_n)) \mid (E_S \cup E_{\text{conn}_1} \cup \cdots \cup E_{\text{conn}_k})$$

Variable-Structure Systems

Variable-structure systems

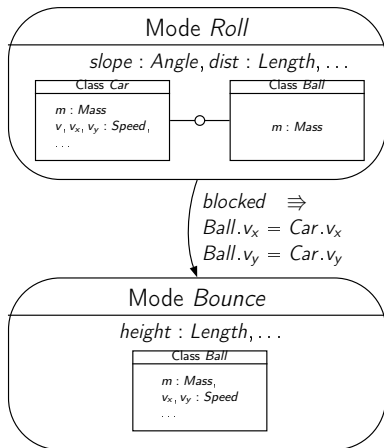
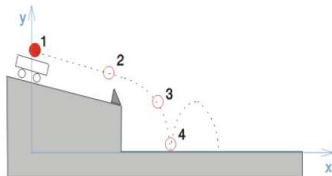
Base case: components with modes



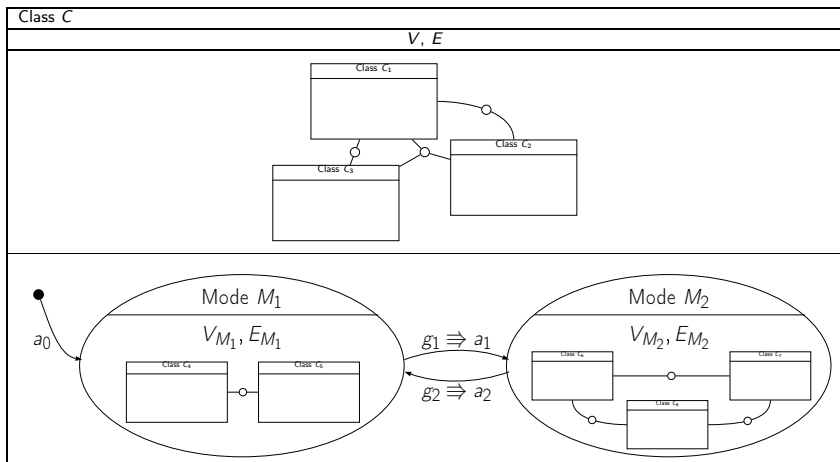
Variable-Structure Systems

Changing Topology

Modes can have different components

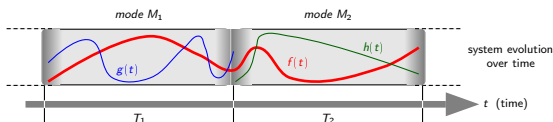


Variable-Structure Systems

The Full Picture of Dynamism

Variable-Structure Systems

Modes



Consider component K of class $C = (V, E, S, D)$:

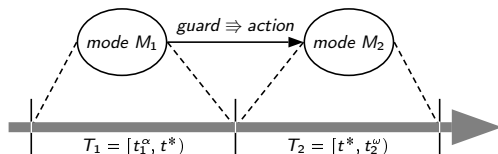
- Component lifetime $T_K = [t^\alpha, t^\omega)$ with $t^\alpha < t^\omega$
- Mode lifetime $T_{M_i} = [t_{M_i}^\alpha, t_{M_i}^\omega)$ (see next slide)
- Semantics during mode M_i : $C_{fix} \otimes S_{M_i} \mid E_{conn}$

Issue: global variables may exhibit different behaviors in different modes

Example: $M_1 : \begin{array}{l} x_1 = 1 \\ x = x_1 \end{array} \quad M_2 : \begin{array}{l} x_2 = 2 \\ x = x_2 \end{array}$

Variable-Structure Systems

Transitions



Transition point t^* is defined by

$t^* = \text{smallest } t, t_1^\alpha < t$ such that

$$guard(t) = true$$

$$\forall \tau, t_1^\alpha < \tau < t. guard(\tau) = false$$

Constraint: $t_1^\alpha < t^* < t_2^\omega$ (modes shall not degenerate to zero length)

Problematic issues to be considered in language design:

- self loops
- conflicting guards (nondeterminism vs. error)

② Simulation Semantics

Numerical solver issues:

- Discretization

ideal functions over \mathbb{R}

$$f : \mathbb{R} \rightarrow \dots$$

are replaced by discrete sampling times

$$\hat{f} : \mathbb{T} \rightarrow \dots$$

where $\mathbb{T} = \{t_i\}$ with $t_{i+1} = t_i + h_i$ for step sizes h_i .

- Precision

- *rounding errors* (limited number size)
- *discretization errors* ($\tilde{x}_{solver} \approx x(t)$)
- *approximation errors* (e.g. Newton algorithm)
- *modeling errors* (parameters, input values)
- *event detection* (e.g. zero crossing)

Solver Semantics

Uncertainty (work in progress)

- Simplest approach is **interval-based**: $x \rightsquigarrow \tilde{x} = x \pm \omega$
- Semantics is defined relative to notion of "**validity**"; hence
 - new concept for validity: $(\tilde{A} \models E)$
 - new concept for models: $\widetilde{Mod}(S) = \{ \tilde{A} \mid \tilde{A} \models S \}$
 - most other constructs (composition etc.) remain unchanged, since $(\dots \otimes \dots \mid \dots)$ is defined relative to validity
- Critical issue: **guards**
 - \tilde{t}^* is defined by $\widetilde{guard}(\tilde{t}^*) = true$
 $\widetilde{guard}(\tau) = false$ for all τ with $\tilde{t}_1^\alpha < \tau < \tilde{t}^*$
 - \rightsquigarrow interval $\tilde{T}_1 = [\tilde{t}_1^\alpha, \tilde{t}^*)$ is blurred
 - \rightsquigarrow **computation traces can be changed** (w.r.t. ideal semantics)
 - \rightsquigarrow analysis techniques are field for intensive research in Numerics

Conclusion

Goal: Compositionality of semantics

- supports variable-structure systems
- supports separate compilation

Goal: Modelica targeted to engineers \rightsquigarrow semantics as well

- semantics streamlined for Modelica
 - \rightsquigarrow no embedded DSL (such as Hydra/Haskell)
- as simple and understandable as possible
 - \rightsquigarrow no large mathematical framework (such as in CIF or Ptolemy)

Goal: Separation of concerns \rightsquigarrow ideal vs. solver semantics

- clear description of modeling principles
- clear description of solver-based constraints
- adaptability to various solvers and solver technologies