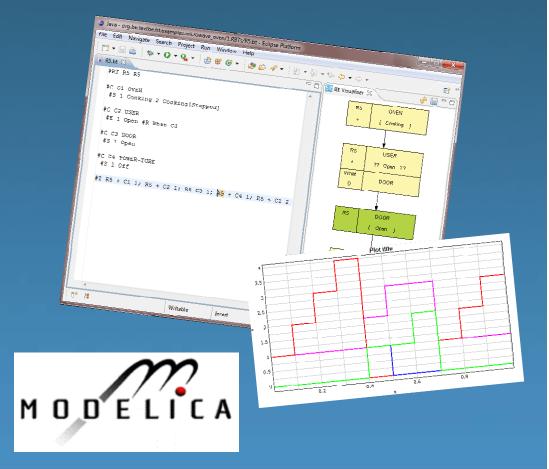
Comodeling Revisited: Execution of Behavior Trees in Modelica

Toby Myers Wladimir Schamai Peter Fritzson





BEHAVIOR

ENGINEERING



Outline

- Introduction to Behavior Engineering
- Comodeling Revisited
- Behavior Trees in Modelica
- Complementing Comodeling with vVDR





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The Software-Hardware Integration Problem Cyber-Physical System Modeling

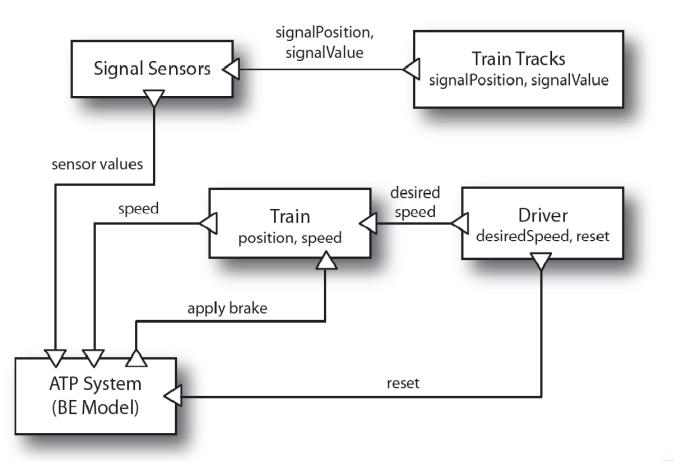
- At the **early stages** of system development, many **decisions** must be made about how the **system will be realised** as a combination of Software and Hardware
- Requirements of the system at these early stages lack quantified and temporal information so it is hard to make an informed decision
- **Changing** the partioning of software / hardware or how they interact later in development can be **time-consuming and costly**
- There is a potential for errors and incompatibility to be introduced as software/hardware specifications are created independently





Example: Model of Automated Train Protection System

An ATP System monitors train position and speed, and may apply brakes if the driver does not react in time







Introduction to Behavior Engineering

Behavior Engineering for Requirements Analysis

- 5 large-scale industry projects
 - In Defence, Transportation, Banking and Finance
 - Between 800-1250 requirements
- All previously reviewed with respective organisations' internal review processes
- Defect detection rate approximately 2 to 3 times that of traditional ad-hoc, checklist-based, and scenario-based reading techniques reported in Porter, 1998.

Requirements Evaluation Using Behavior Trees

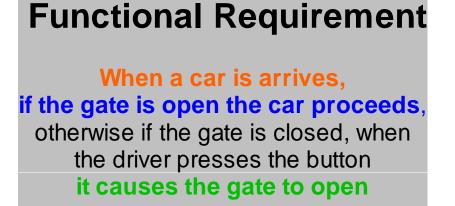
Findings from Industry

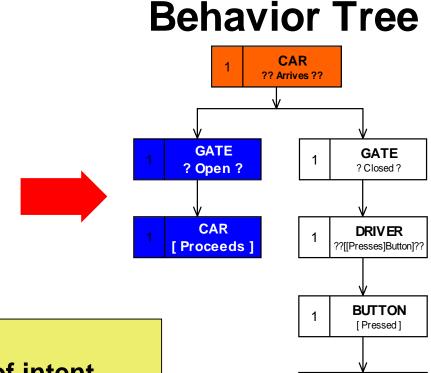
Daniel Powell http://aswec07.cs.latrobe.edu.au/5.zip





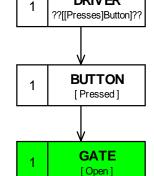
Formalization - Requirements Translation





Formalization

- clarification and preservation of intent
- strict use of original vocabulary
- removes ambiguity, aliases, etc
- aids stakeholder validation, understanding
- approaches repeatability







A Brief Introduction to Behavior Engineering

• Behavior Engineering is a methodology with a tightly interlinked language and process

Behavior Modeling Process	Behavior Modeling Language (BML)		
(BMP)	Behavior Trees (BT)	Composition Trees (CT)	
Requirements Translation	Requirement Behavior Trees (RBTs)	Requirement Composition Tree (RCT)	
Requirements Integration	Integrated Behavior Tree (IBT)	Integrated Composition Tree (ICT)	
System Specification	Model Behavior Tree (MBT)	Model Composition Tree (MCT)	
System Design	Design Behavior Tree (DBT)	Design Composition Tree (DCT)	





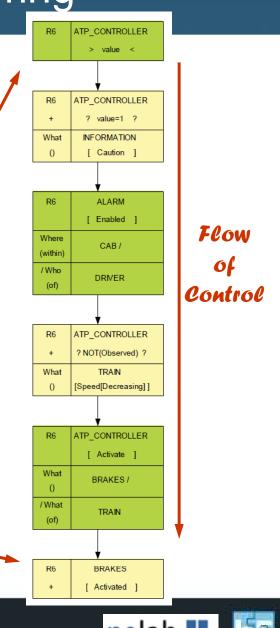
Introduction to Behavior Engineering

How to translate from a Requirement in Natural Language to an RBT

R6. If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's cab. Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.

The Tag traces these Behavior Tree nodes back to Requirement 6.

A '+' and a yellow color denote the behavior is implied by the requirements







A Brief Introduction to Behavior Engineering

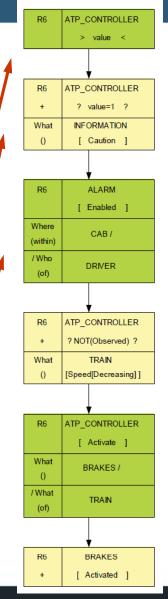
How to translate from a Requirement in Natural Language to an RBT

R6. If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's cab. Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.

ATP Controller receives a value from another component

Check if the value is a cantion signal

If it is, enable the Alarm. To maintain the intent of the original requirement, use a relation to show the Alarm is enabled in the Driver's Cab.







A Brief Introduction to Behavior Engineering

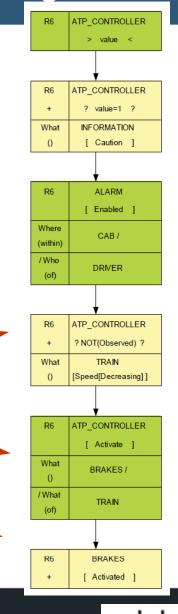
How to translate from a Requirement in Natural Language to an RBT

R6. If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's cab. Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.

It is implied the ATP Controller must observe whether the Train's speed is decreasing.

If the Train isn't decreasing in speed, the ATP Controller activates the Braking System of the Train.

... Which results in the Braking System being Activated -







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Comodeling of Cyber-Physical Systems Revisited

- At the **early stages** of system development, many **decisions** must be made about how the **system will be realised** as a combination of Software and Hardware
- Requirements of the system at these early stages lack quantified and temporal information so it is hard to make an informed decision
- **Changing** the partioning of software / hardware or how they interact later in development can be **time-consuming and costly**
- There is a potential for errors and incompatibility to be introduced as software/hardware specifications are created independently





Requirement	Description
RI	The ATP system is located on board the train. It involves a central controller and five boundary subsystems that manage the sensors, speedometer, brakes, alarm and a reset mechanism.
R2	The sensors are attached to the side of the train and detect information on the approach to track-side signals, i.e. they detect what the signal is displaying to the train driver.
R3	In order to reduce the effects of component failure three sensors are used. Each sensor generates a value in the range 0 to 3, where 0, 1 and 2 denote the danger, caution, and proceed signals respectively. The fourth sensor value, i.e. 3, is generated if an undefined signal is detected, e.g. may correspond to noise between the signal and the sensor.
R4	The sensor value returned to the ATP controller is calculated as the majority of the three sensor readings. If there does not exist a majority then an undefined value is returned to the ATP controller.
R5	If a proceed signal is returned to the ATP controller then no action is taken with respect to the train's brakes.
R6	If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's cab. Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.
R7	In the case of a danger signal being returned to the ATP controller, the braking system is immediately activated and the alarm is enabled. Once enabled, the alarm is disabled if a proceed signal is subsequently returned to the ATP controller.
R8	Note that if the braking system is activated then the ATP controller ignores all sensor input until the system has been reset.
R9	If enabled, the reset mechanism deactivates the train's brakes and disables the alarm.

 Table 1. Requirements of the ATP system





Interaction with Sensors ...

Requirement	Description	
R1	The ATP system is located on board the train. It involves a central controller and five boundary subsystems that manage the sensors, speedometer, brakes, alarm and a reset mechanism.	
R2	The sensors are attached to the side of the train and detect information on the approach to track-side signals, i.e. they detect what the signal is displaying to the train driver.	How often
R3	In order to reduce the effects of component failure three sensors are used. Each sensor generates a value in the range 0 to 3, where 0, 1 and 2 denote the danger, caution, and proceed signals respectively. The fourth sensor value, i.e. 3, is generated if an undefined signal is detected, e.g. may correspond to noise between the signal and the sensor.	How often does this need to be checked?
R4	The sensor value returned to the ATP controller is calculated as the majority of the three sensor readings. If there does not exist a majority then an undefined value is returned to the ATP controller.	
R5	If a proceed signal is returned to the ATP controller then no action is taken with respect to the train't brakes.	
R6	If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's each Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.	
R7	In the case of a danger signal being returned to the ATP controller, the braking system is immediately activated and the alarm is enabled. Once enabled, the alarm is disabled if a proceed signal is subsequently returned to the ATP controller.	Decreasing by how much?
R8	Note that if the braking system is activated then the ATP controller ignores all sensor input until the system has been reset.	by how much?
R9	If enabled, the reset mechanism deactivates the train's brakes and disables the alarm.	

Table 1. Requirements of the ATP system



Interaction with Actuators ...

Requirement	Description	=
RI	The ATP system is located on board the train. It involves a central controller and five boundary subsystems that manage the sensors, speedometer, brakes, alarm and a reset mechanism.	=
R2	The sensors are attached to the side of the train and detect information on the approach to track-side signals, i.e. they detect what the signal is displaying to the train driver.	-
R3	In order to reduce the effects of component failure three sensors are used. Each sensor generates a value in the range 0 to 3, where 0, 1 and 2 denote the danger, caution, and proceed signals respectively. The fourth sensor value, i.e. 3, is generated if an undefined signal is detected, e.g. may correspond to noise between the signal and the sensor.	-
R4	The sensor value returned to the ATP controller is calculated as the majority of the three sensor readings. If there does not exist a majority then an undefined value is returned to the ATP controller.	-
R5	If a proceed signal is returned to the ATP controller then no action is taken with respect to the train's brakes.	what response
R6	If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's cab. Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.	time is realistically acceptable?
R7	In the case of a danger signal being returned to the ATP controller, the braking system is immediately activated and the alarm is enabled. Once enabled, the alarm is disabled if a proceed signal is subsequently returned to the ATP controller.	
R8	Note that if the braking system is activated then the ATP controller ignores all sensor input until the system has been reset.	-
R9	If enabled, the reset mechanism deactivates the train's brakes and disables the alarm.	

Table 1. Requirements of the ATP system





Software / Hardware Partitioning ...

Requirement	Description	
R1	The ATP system is located on board the train. It involves a central controller and five boundary subsystems that manage the sensors, speedometer, brakes, alarm and a reset mechanism.	Dankansala
R2	The sensors are attached to the side of the train and detect information on the approach to track-side signals, i.e. they detect what the signal is displaying to the train driver.	Perform in Software
R3	In order to reduce the effects of component failure three sensors are used. Each sensor generates a value in the range 0 to 3, where 0, 1 and 2 denote the danger, caution, and proceed signals respectively. The fourth sensor value, i.e. 3, is generated if an undefined signal is detected, e.g. may correspond to noise between the signal and the sensor.	or Hardware?
R4	The sensor value returned to the ATP controller is calculated as the majority of the three sensor readings. If there does not exist a majority then an undefined value is returned to the ATP controller.	
R5	If a proceed signal is returned to the ATP controller then no action is taken with respect to the train's brakes.	
R6	If a caution signal is returned to the ATP controller then the alarm is enabled within the driver's cab. Furthermore, once the alarm has been enabled, if the speed of the train is not observed to be decreasing then the ATP controller activates the train's braking system.	
R7	In the case of a danger signal being returned to the ATP controller, the braking system is immediately activated and the alarm is enabled. Once enabled, the alarm is disabled if a proceed signal is subsequently returned to the ATP controller.	
R8	Note that if the braking system is activated then the ATP controller ignores all sensor input until the system has been reset.	
R9	If enabled, the reset mechanism deactivates the train's brakes and disables the alarm.	

Table 1. Requirements of the ATP system





The Environment in which the system will exist ...

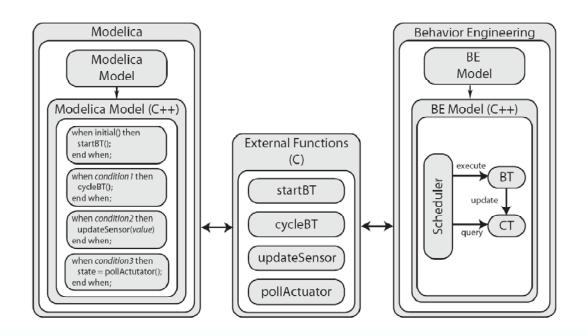
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d, the reset mechanism deactivates the train's brakes and divables the alarm.	
A as t t	TP controller activates the train's braking system. e of a danger signal being returned to the ATP controller, the braking system is immediately and the alarm is enabled. Once enabled, the alarm is disabled if a proceed signal is subsequently to the ATP controller. if the braking system is activated then the ATP controller ignores all sensor input until the system reset.







- Previous implementation of Comodeling
 - BE and Modelica models executed separately
 - Integrated using Modelica external functions







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- Why execute Behavior Trees in Modelica?
 - 1. Comodeling is easier to apply as the comodel is captured solely in Modelica
 - A Behavior Tree can directly affect the acausal equations used to model the hardware and environmental components
 - Modelica is used as an **action language** for the whole integrated Co-model
 - 2. Portable models the ability to execute Behavior Trees is available to a wide audience, several tools
 - Comodeling can be used with other complementary approaches such as the virtual verification of system designs against system requirements





- Representing Behavior Trees in Modelica
 - Basic Nodes
 - State Realisation, Selection, Guard, Input, Output
 - Branching & Composition
 - Sequential Composition, Atomic Composition, Parallel Branching, Alternative Branching
 - Operators
 - Reference, Reversion, Branch-Kill, Synchronisation





- State Realisation
 - A state realisation updates the *state* variable of the associated component to the enumerated value of the behavior of the BT node.

c.state := Integer(c.states.s);





- Selection
 - A selection performs an equality check on the state variable of the associated component, comparing it to the enumerated value of the behavior of the BT node.
 Depending on the result of this equality check, the flow of control either continues or is terminated.

```
if c.state == c.state_s then
    ... // Continue flow of control
else
    ... // Terminate branch
end if;
```





- Guard
 - A guard is similar to a selection, with the exception that the else branch is not included to ensure that the guard is continually re-evaluated until true.

```
if c.state == c.state_s then
    ... // Continue flow of control
end if;
```





- Input
 - Inputs and outputs are implemented as boolean variables. When the variable is true, the input is active. Events last for one cycle, to ensure that if an internal output is active in one branch, it can be received by an internal output in another branch. Inputs are represented with an equality check that is true if the associated event becomes active.

if e2 then
 ... // Continue flow of control
end if;





- Sequential Composition
 - Updates the value of the branch variable

if branch1 == 1 then ... // Node behavior branch1 := 2; elseif ...





- Parallel Branching
 - Clears the current branch value and sets the branch value of the child branches to their first node.

if branch1 == 1 then ... // Node behavior branch1 := 0; branch2 := 1; branch3 := 1; elseif ...





- Alternative Branching
 - As per parallel branch, but when the first node of any of the child branches is activated all the sibling branches are terminated.

if branch2 == 1 then
 ... // Node behavior
 branch2 := 2;
 branch3 := 0;
end if;





- Atomic Composition
 - Adds further constraint to all sibling branches of atomic composed nodes that flow of control cannot continue if the branch values of the atomic composed nodes are active.

if branch3 == 1 and not(branch2==1 or branch2==2) then ... // Node behavior branch3 := 2;





- Reference
 - Clears the current branch value and sets the branch value of the destination node.

if branch3 == 1 then
 branch3 := 0;
 branch2 := 2;





- Reversion
 - Clears the current branch value and all sibling parallel branches and sets the branch value of the destination node.

if branch1 == 3 then
 branch2 := 0;
 branch3 := 0;
 branch1 := 1;





- Branch-Kill
 - Clears the branch value of the destination node and the branch value of any of its descendants.

if branch1 == 2 then
branch3 := 0;
... // Continue flow of control





- Synchronisation
 - One synchronisation node checks when the branch value of all nodes is set correctly and sets a boolean variable to true. All other synchronisation nodes wait until this Boolean variable is true.

if branch3 == 2 and branch2 == 3 then
 sync1 := true;
 ... // Node behavior
 ... // Continue flow of control
if branch2 == 3 and sync1 then
 ... // Continue flow of control





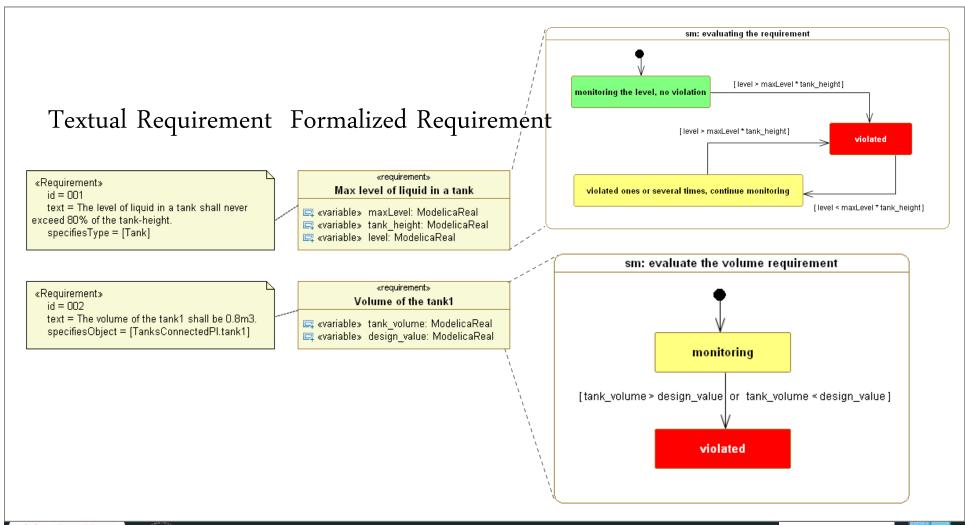
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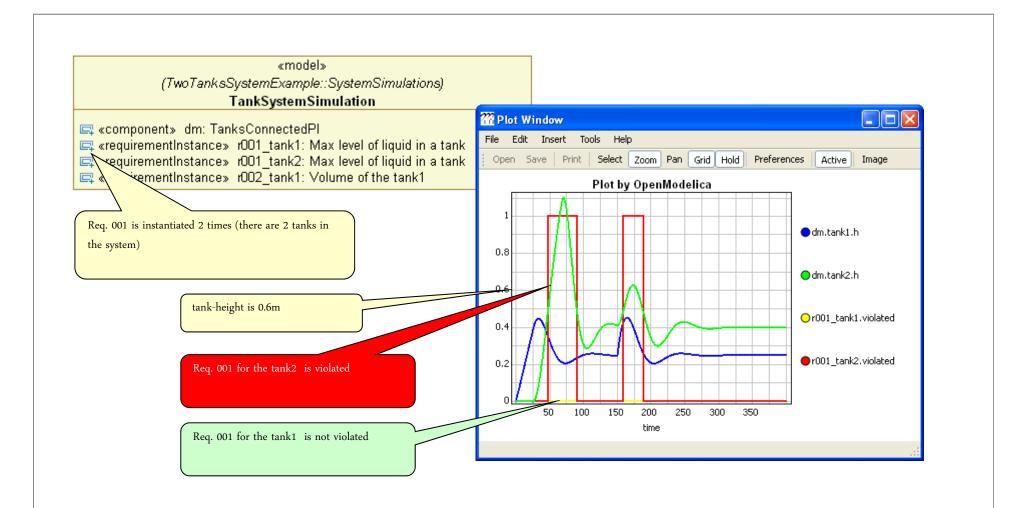
vVDR – Virtual Verification of Design Requirements ModelicaML Example: System Requirements



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vVDR – Virtual Verification of Design Requirements in ModelicaML Simulation and Requirements Check







Complementing Comodeling with vVDR

- vVDR and BE provide differing benefits
 - vVDR can verify different system design alternatives against the same set of requirements or drive verifications driven by different test scenarios.
 - BE provides a means to ensure the consistency of a specification.
- Their combination could leverage the benefits of both approaches





Complementing Comodeling with vVDR

- vVDR is a method for virtual Verification of system Design alternatives against system Requirements
- In vVDR each requirement is formalised as a model to evaluate violation and fulfilment criteria
- These requirement models and a system design are instantiated and bound into a test model which is translated into Modelica





Complementing Comodeling with vVDR

- Two possible applications
 - Integrate BE with vVDR by using a model behavior tree as the source for generation of both vVDR requirements violation monitors and test cases
 - 2. Augment the existing comodeling approach with vVDR. vVDR provides monitors that evaluate the performance of different comodels to find the best candidate that fulfils a set of criteria





Implementation Aspects

- Use declarative (equations) or imperative (algorithms) constructs?
- Create a Modelica library or create a code generator?
- In both cases, Behavior Trees in Modelica, and ModelicaML state charts, a code generator to algorithmic Modelica was implemented







• Any Questions?

- vVDR support in ModelicaML available from www.openmodelica.org
- A little plug ...



- <u>www.tjmyersconsulting.com</u>
- Specialise in practicing Behavior Engineering



