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Towards a Computer Algebra System with Automatic Differentiation for use with object-oriented modelling languages

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OPTEC - Optimization in Engineering

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- Interdiciplinary: Mech.Eng. + Elec.Eng. + Civ.Eng. + Comp.Sc.
- Katholieke Universiteit Leuven, Belgium
- 2005-2010, phase II 2010-2017



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Myself

- M.Sc. Engineering Physics/Mathematics from Chalmers, Gothenburg
- PhD student since Oct 2008 for Prof. Moritz Diehl
- Topic: Modelling and Derivative Generation for Dynamic Optimization and Application to Large Scale Interconnected DAE Systems
- Application: Solar thermal power plant





Dynamic optimization problem

We consider dynamic optimization problems of the form (can be generalized further):

 $\begin{array}{lll} \underset{x(\cdot), z(\cdot), u(\cdot), p}{\text{minimize:}} & \int_{t=0}^{T} L(x, u, z, p, t) \, dt + E(x(T)) \\ & \\ f(\dot{x}(t), x(t), z(t), u(t), p, t) & = 0 & t \in [0, T] \\ & \\ \text{subject to:} & h(x(t), z(t), u(t), p, t) & \leq 0 & t \in [0, T] \\ & \\ f(x(0), x(T), p) & = 0 \end{array}$

 $x:\mathbb{R}_+\to\mathbb{R}^{N_x}$ differential states, $z:\mathbb{R}_+\to\mathbb{R}^{N_z}$ algebraic states, $u:\mathbb{R}_+\to\mathbb{R}^{N_u}$ control, $\rho\in\mathbb{R}^{N_p}$ free parameters



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Solution

Dynamic programming / Hamilton-Jacobi-Bellman equation – for very small problems

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• Multiple-shooting: Parametrize state at some times and use single shooting in between

Good reference: L. Biegler Nonlinear Programming, SIAM 2010

Direct Multiple Shooting (Bock, 1984)

- Subdivide time horizon: $0 = t_0 \leq \ldots \leq T_N$
- Parametrize control: $u(t) = u_i$, $t \in [t_i, t_{i+1}]$
- Parametrize state: s_{x,i} = x(t_i)



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- Nonlinear Program (NLP):

$$\begin{array}{ll} \text{minimize:} & \sum_{s_{x,i}, u_i, p}^{N-1} \sum_{i=0}^{N-1} L_i(s_{x,i}, u_i, p) + E(s_{x,N}) \end{array}$$

subject to:

$$\begin{array}{lll} s_{x,i+1} &= F_i(s_{x,i},u_i,p), & \forall \\ 0 &\geq h_i(s_{x,i},u_i,p), & \forall \\ 0 &= r(s_{x,0},s_{x,N},p) \end{array}$$

F_i: Call to an DAE integrator



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- Solve with e.g. structure-exploiting SQP method
- Software: ACADO Toolkit, MUSCOD-II



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Notes

• The problem formulation can be generalized (e.g. free end time, multiple model stages, hybrid)

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 Large scale NLP solvers require derivative information, at least first order

Automatic differentiation

Automatic differentiation

Automatic differentiation^a, AD, is able to cheaply and accurately providing derivative evaluation of a function f = f(x) by applying the **chain rule to the algorithm**. Two "modes":

- Forward mode: $\frac{\partial f}{\partial x} r$
- Reverse (adjoint) mode: r^T <u>∂f</u>/_{∂x}

^aSee e.g. Griewank & Walther: Evaluating Derivatives, 2008



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Implementations

Operator overloading (OO)

- Idea: Use operator overloading in e.g. C++, to record calculations
- Easy to implement, expecially in forward mode, needs only a C++- compiler
- Disadvantage: Effective implementation is based on template meta programming

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- Tools: ADOL-C, CppAD, ...
- Source code transformation (SCT)
 - Implements AD inside a compiler (think GCC)
 - Advantage: Efficient code
 - Disadvantage: Hard to implement, less mature than OO
 - Tool: OpenAD

Can be applied to "black-box" C or Fortran functions!

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Motivation

Observations

- To apply an existing AD tool, we first need to generate C-code
- The AD tool will parses the code to obtain the graph we already had!
- More sensible approach: Apply AD to the graph directly

Benifits:

- No compiler in the loop
- No information losses (e.g. for systems with switches)
- Convex reformulation (cf. CVX)
- Structure exploatation



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Structure exploatation: The Lifted Newton method

- Albersmeyer & Diehl, SIAM 2010
- "Lifts" non-linear root-finding problems to a higher dimension
- Speeds up convergence, increases region of attraction
- Requires that the functions are given as an algorithm
- Example: A single-shooting algorithm can be lifted to a multiple-shooting algorithm with condensing

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What is CasADi?

A minimalistic Computer Algebra System (CAS) written in self-contained C++

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Tailored for high speed, especially during numerical evaluation



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- Tailored for high speed, especially during numerical evaluation
- Contains interfaces to Sundials (IDAS,CVodes), IPOPT, ACADO Toolkit, JModelica
- Open-source project: first public release autumn 2010 (www.casadi.org)
- Permissive licence, LGPL

Structure of CasADi

Two graph representations used in conjunction

	Scalar expression, SX "maximum speed"	Matrix expression, MX "maximum generality"
number of arguments	two	arbitrary
nodes	scalar-valued	vector-valued
operations	built-in	all (e.g. calls to FX)
branching/jumps	no	yes
parallelization	no	yes
syntax	double	double, ublas::matrix

Dynamically created functions

Function expression, FX

- $[y_1,\ldots,y_n] = f(x_1,\ldots,x_m), y_i \in \mathbb{R}^{r_i}, x_j \in \mathbb{R}^{q_j}$
- Polymorphic class, derived classes:
 - Function given by a graph of SX nodes
 - Function given by a graph of MX nodes
 - External function (e.g. from DLL)
 - Integrator and "Simulator"

Observation

- SX \sim DAG in AD-tools
- $\bullet \ \mathsf{MX} \sim \mathsf{DAG} \text{ in Modelica}$



The Integrator class

- An *integrator* is considered to be a function: $x(t_f) = F(t_0, t_f, x(t_0), p)$
- Currently one explicit integrator (CVodes) and one implicit integrator (IDAS)
- AD forward/reverse corresponds to forward/adjoint sensitivities



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The Simulator class

• Evaluates an output function, h(t, x, p) in a set of time points, $[t_1, \ldots, t_n]$, using an arbitrary integrator

• Also considered a function: $[y_1, \ldots, y_n, x(t_f)] = G([t_0, t_f, x(t_0), p, [t_1, \ldots, t_n])$ with $y_i := h(t_i, x(t_i), p)$



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Relation to dynamic optimization

The integrator and simulator classes used e.g. in shooting-methods

Determinant calculation

- AD speed benchmark (ADOL-C, CppAD)
- $f(X) = |X|, \quad X \in \mathbb{R}^{N \times N}$ by minor expansion
- Complexity exponential in N
- Adjoint derivatives
- Intel Core Duo 2.4 GHz, 4 GB RAM, 3072 KB L2 cache, 128 KB L1 cache



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Results

- Outperforms ADOL-C, keeps up with CppAD up to 8-by-8 (\approx 100.000 operations)
- 20 ns per operation (= speed of cache)
- Optimized c-code (generated), not much faster





Minimal fuel rocket flight

minimize: u, s, v, m	-m(T)	
	s = v $\dot{v} = (u - \alpha v^2)/m$	
	$\dot{m} = -\beta u^2$	(1)
subject to:	s(0) = 0, s(T) = 10 v(0) = 0, v(T) = 0	
	m(0) = 0, v(T) = 0 m(0) = 1, T = 10	
	$-10 \leq u \leq 10$	

Euler forward integrator

```
SX s_0("s_0"), v_0("v_0"), m_0("m_0");
std::vectorSX> x0 = {s_0,v_0,m_0};
SX u("u");
SX s = s_0, v = v_0, m = m_0;
double dt = 10.0/1000;
for(int j=0; j<1000; ++j){
    s += dt*v;
    v += dt / m * (u - alpha * v*v);
    m += -dt * beta*u*u;
}
std::vector<SX> x = {s,v,m};
FX integrator = SXFunction({x0,u},x);
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Single shooting

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std::vector<double> X0 = {0,0,1}; // X at t=0
MX X = X0; // state vector
MX U("U",1000); // control vector
for(int k=0; k<1000; ++k){
    X = integrator.evaluate({X,U[k]});
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MX s_T = X[0], v_T = X[1], m_T = X[2];
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1000 controls intervals, 1000 steps per interval

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- 1000 controls intervals, 1000 steps per interval
- Build graph for integrating over a single interval
- This graph defines the function "integrator"
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- Build graph for integrating over a single interval
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- This graph defines objective and constraint funcs.
- Pass to NLP solver (IPOPT)

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- Convergence after 11 iteration, 10.4 s.

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}
MX s_T = X[0], v_T = X[1], m_T = X[2];
...</pre>
```

OPTEC

Van-der-pol oscillator

- From JModelica example collection
- Export OCP as Modelica/Optimica XML
- Parse the XML code in CasADi, reconstruct OCP
- Solve with ACADO Toolkit
 - Open-source dynamic optimization software from OPTEC
 - Houska & Ferreau 2008-present
 - www.acadotoolkit.org
 - Spatial discretization with multiple-shooting
 - Non-linear program (NLP) solved with Sequential Quadratic Programming (SQP)
 - Limited memory Hessian approximation
 - Initialized with u = 0 for all t

Convergence after 26 iterations







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- First public release of CasADi (autumn 2010)
- Fully integrate with JModelica, support all features
- Applications: solar power, combined cycle, hydropower valley
- Large-scale (distributed, parallel) SQP (with C. Savorgnan, A. Kozma)



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Further research

- Dynamic optimization of large-scale hybrid systems (reformulate as MINLP)
- PDE constrained optimization





Thank you for listening.

