



EADS INNOVATION WORKS

Systems Engineering



Execution of UML State Machines Using Modelica

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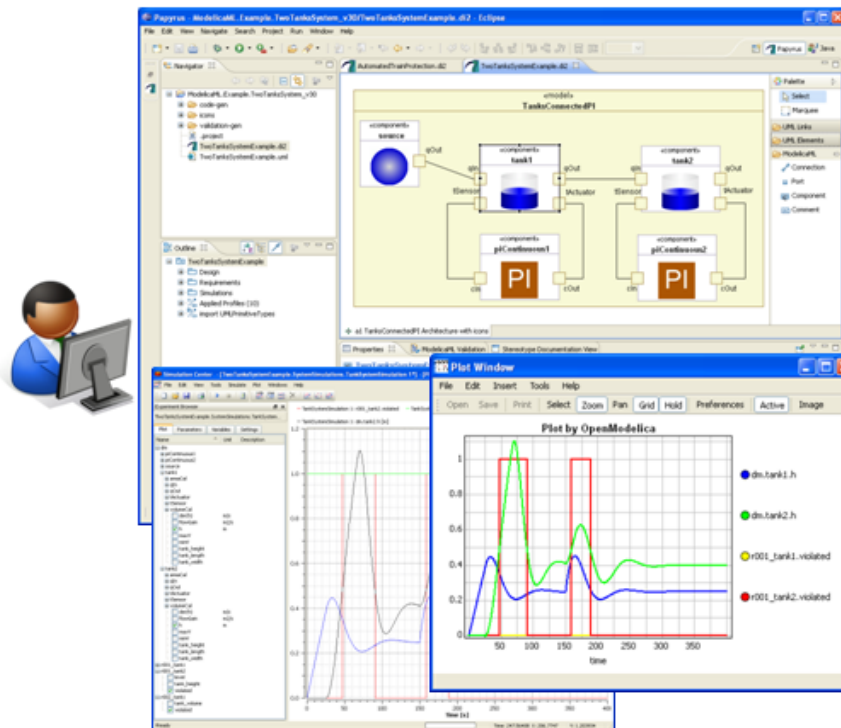
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Introduction: ModelicaML Concept

① System Modeling with ModelicaML



② Modelica Code Generation

```

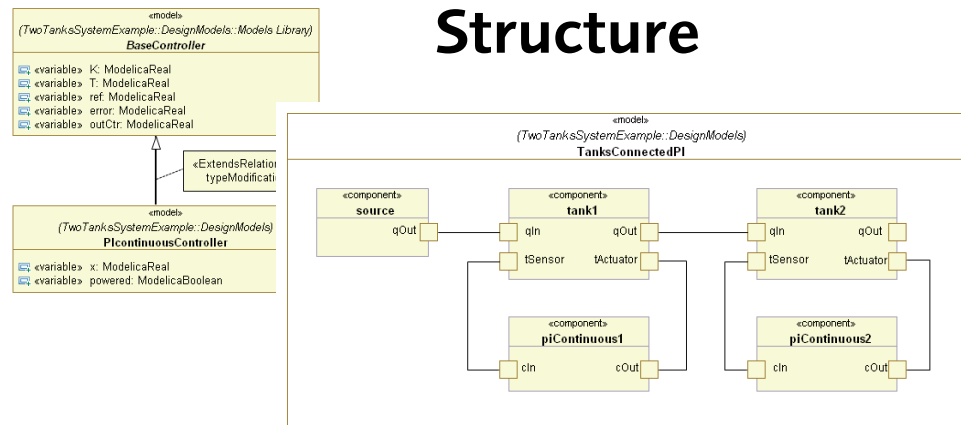
1  within TwoTanks;
2
3  function limitValue
4
5  input Real pRin;
6  input Real pRin;
7  output Real pLin;
8  input Real p;
9
10 algorithm
11 // code generated from the Activity "Algorithm (diagram)"
12 // Activity name: "Algorithm (diagram)"
13 if p < pRin then
14   pLin := pRin; // Operation name: "pLin := pRin"
15 else if p > pRin then
16   pLin := p;
17 end if;
18
19 model Task
20
21 //TaskTypeExample.DesignNode.LiquidSource source(FlowLevel = 0.02);
22 //TaskTypeExample.DesignNode.LiquidFlow qIn //Connector, flow (m³/s) through input
23 //TaskTypeExample.DesignNode.LiquidFlow qOut //Connector, flow (m³/s) through output
24 parameter Real FlowLevelInit = "qIn" = 0.02;
25 parameter Real qIn = 0; //Liquid flow input value flow
26 parameter Real qOut = 0; //Liquid flow output value flow
27 Real kControl = 0.1, uSet = "uSet" //Task level
28 parameter Real taskLength = 0.1;
29 parameter Real taskWidth = 1;
30 parameter Real taskLength = 1;
31
32 within DesignNode;
33
34 model TaskConnectedPI
35
36 //TaskTypeExample.DesignNode.LiquidSource source(FlowLevel = 0.02);
37 //TaskTypeExample.DesignNode.LiquidFlow qIn //Connector, flow (m³/s) through input
38 //TaskTypeExample.DesignNode.LiquidFlow qOut //Connector, flow (m³/s) through output
39 //TaskTypeExample.DesignNode.PIDController pControl(pControlInit = 0.25);
40 //TaskTypeExample.DesignNode.PIDController pControl(pControlInit = 0.4);
41
42 equation
43 connect(source.qOut, task1.qIn);
44 connect(pControl1.cIn, task1.tSensor);
45 connect(pControl1.qOut, task1.tActuator);
46 connect(task1.qOut, task2.qIn);
47 connect(pControl2.cIn, task2.tSensor);
48 connect(pControl2.qOut, task2.tActuator);
49
50 end TaskConnectedPI;

```

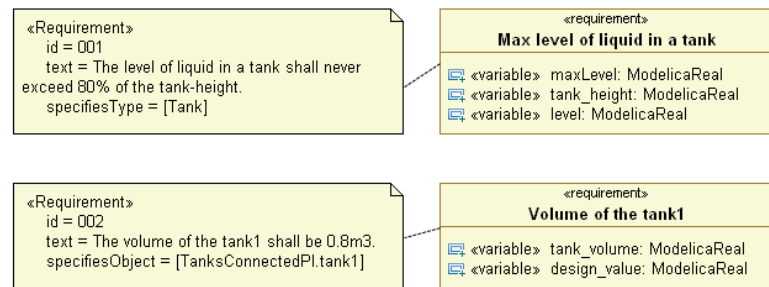
③ System Simulation with Modelica Tools

Introduction: ModelicaML Graphical Notation

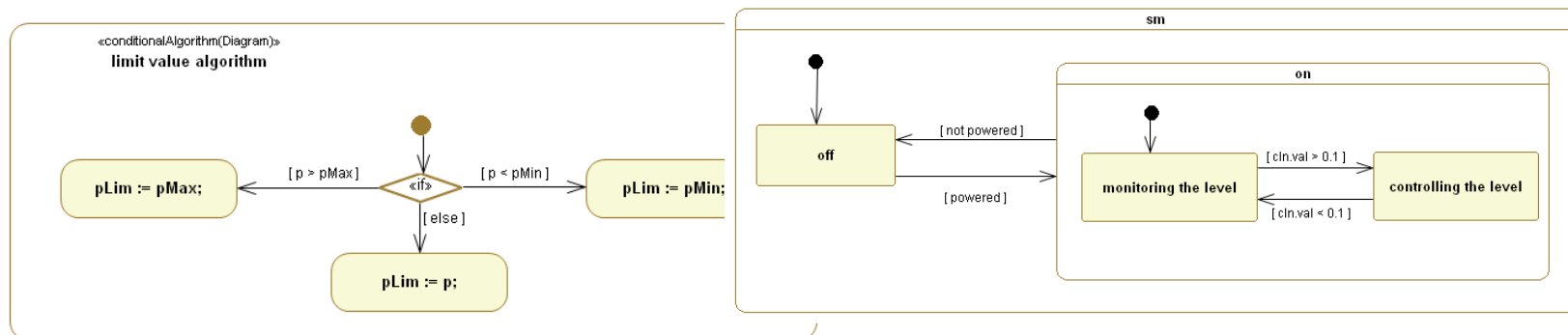
Structure



Requirements



Behavior



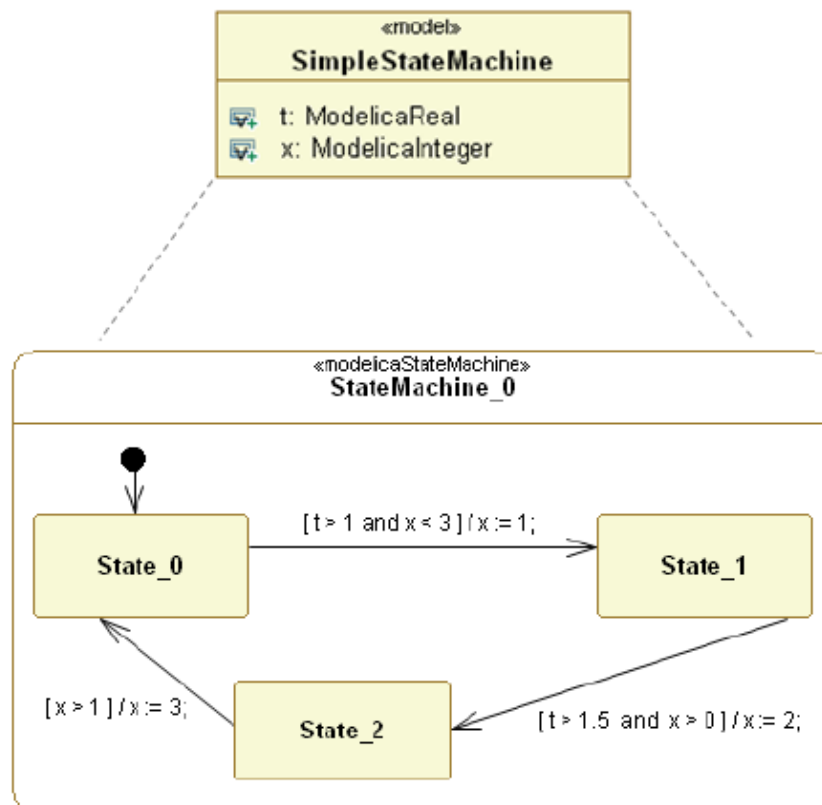
Introduction: Motivation

- UML state machines is a powerful formalism for describing system or components states and modes
- To which extent can UML state machines concepts be supported in ModelicaML?
- How to implement it?
 - Create a Modelica library or implement a code generator?
 - Use algorithmic code or equations?
- This presentation focuses on the resolution of issues identified when translating UML state machines into executable Modelica code

Translation Approach

- Code generator approach was chosen instead library
 - States specification (i.e. in/outgoing transitions, entry, do, exit actions) cannot be predefined as a library component.
- ModelicaML code generator generates algorithmic Modelica code from ModelicaML state machines because:
 - The behavior of a state machine is always causal
 - For inter-level transitions, i.e. transitions which cross states hierarchy borders, the deactivation and activation of states and the execution sequence of associated actions (exit/entry action of states or state transitions effects) has to be performed in an explicitly defined order.

Simple Example (1/2)



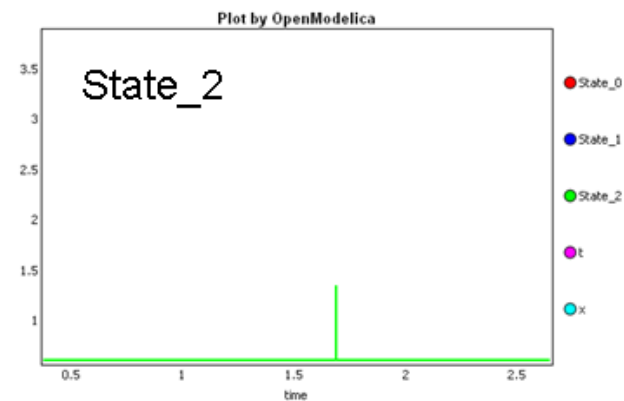
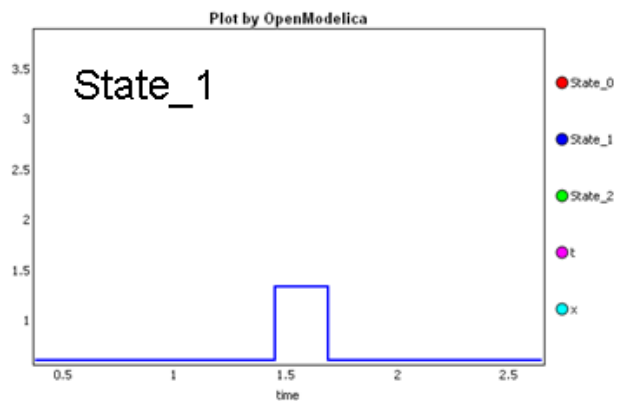
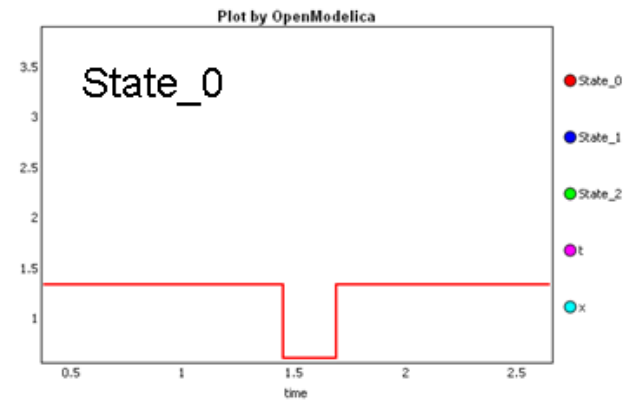
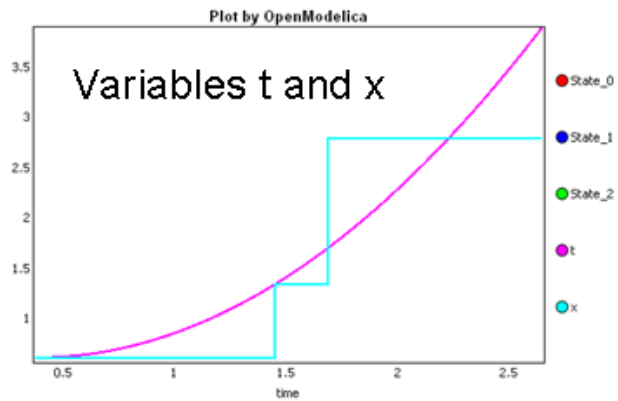
```

model SimpleStateMachine
  Boolean State_0 "State_0 representation";
  Boolean State_1 "State_1 representation";
  Boolean State_2 "State_2 representation";
  Integer x "discrete variable";
  Real t "continuous variable";
  equation //Code for continuous integration of t
    der(t)=time;

  algorithm //Code for StateMachine_0
    when initial() then
      State_0:=true "Activation of the initial state";
    end when;

    //Transition from State_0 to State_1
    if pre(State_0) and t > 1 and x < 3 then
      State_0:=false "Deactivation of state";
      x:=1 "Transistion effect";
      State_1:=true "Activation of state";
    //Transition from State_1 to State_2
    elseif pre(State_1) and t > 1.5 and x > 0 then
      State_1:=false "Deactivation of state";
      x:=2 "Transistion effect";
      State_2:=true "Activation of state";
    //Transition from State_2 to State_0
    elseif pre(State_2) and x > 1 then
      State_2:=false "Deactivation of state";
      x:=3 "Transistion effect";
      State_0:=true "Activation of state";
    end if;
  end SimpleStateMachine;
  
```

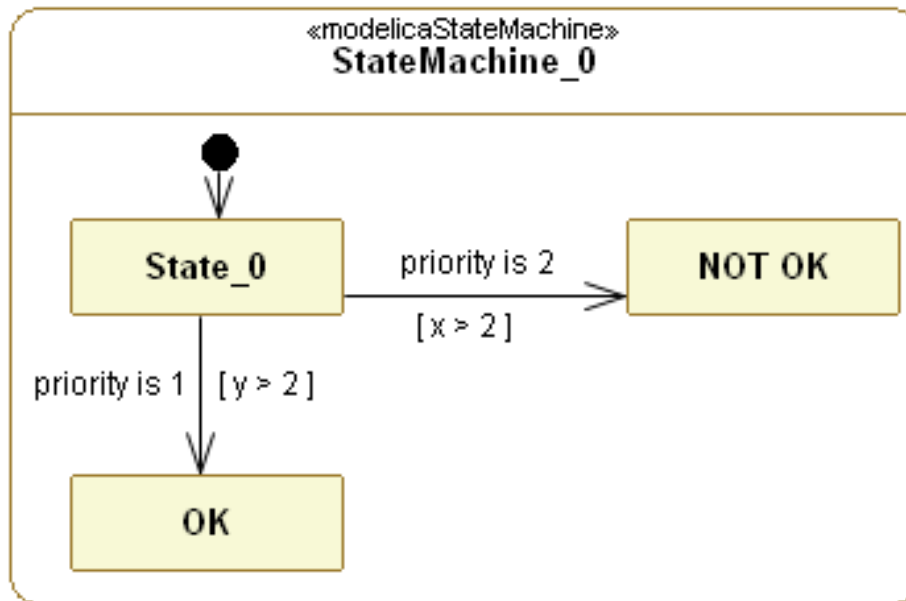
Simple Example (2/2)



UML State Machine Concepts Supported in ModelicaML

- Hierarchical states modeling
 - Composite states
 - Sub-state machines (reusable for multiple states)
- Regions (orthogonal states)
- Pseudo states
 - initial , shallowHistory, join, fork, junction, choice, entryPoint, exitPoint, terminate
- Transition
 - Compound transitions (a transitions set from state to state through pseudo states)
 - Inter-level transition (transitions that cross hierarchy levels)
- Events
 - Change Events, Time Events, Signal Events
- State Actions
 - onEntry, Do, onExit

Issues with Conflicting Transitions



What happens when x and y are greater than 2 at the same time?

In ModelicaML: *Transition which the higher execution priority is taken (i.e. to state OK)*

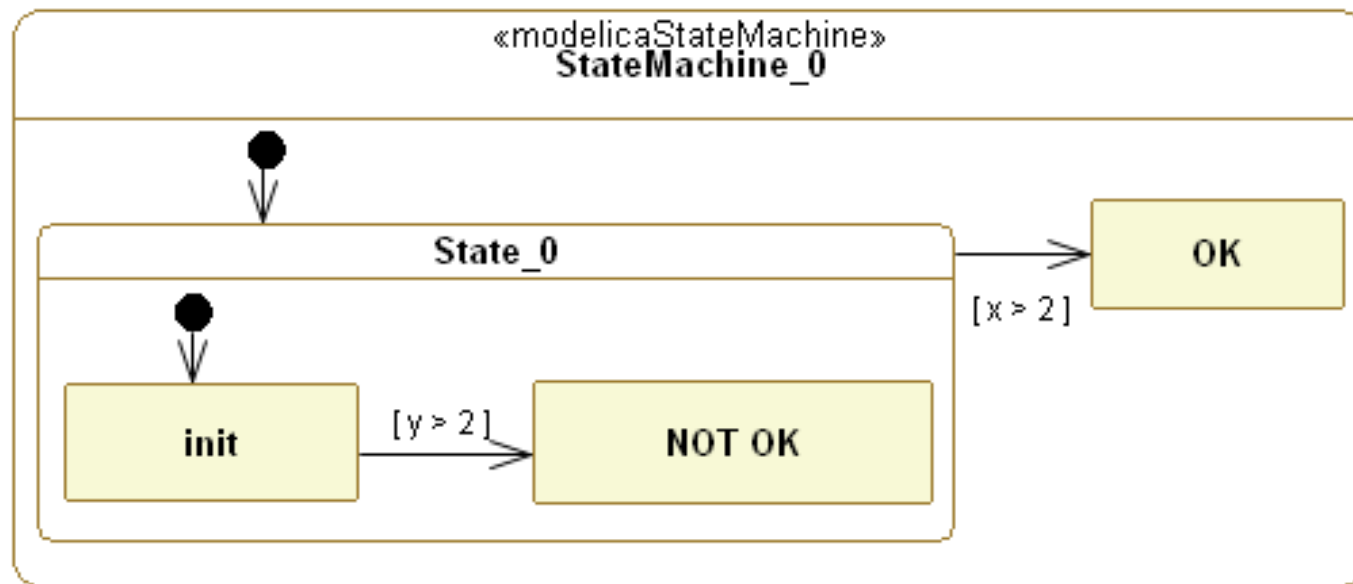
Priority Order

```

1: target-> OK[x > 1]
2: target-> NOT OK[x > 2]
  
```



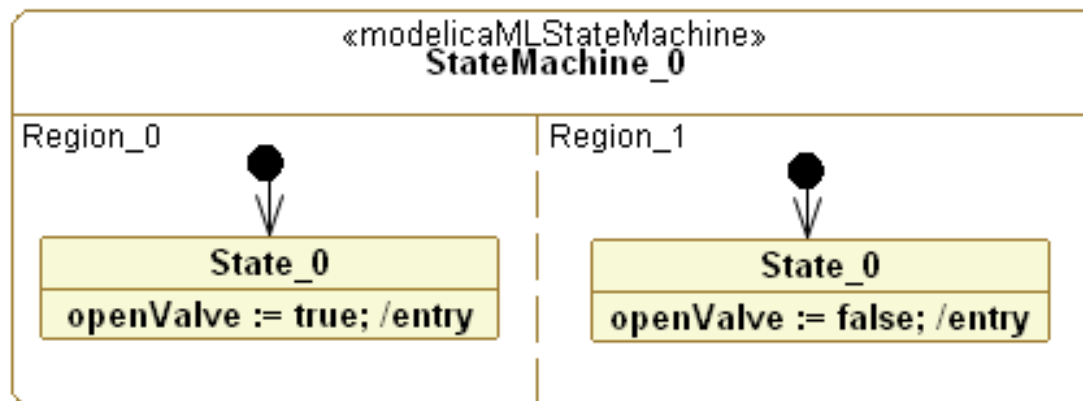
Priority Schema for Conflicting Transitions at Different State Hierarchy Levels



What happens when **x** and **y** are greater than **2** at the same time?

In ModelicaML: Transition to state "OK" is taken because it is at higher hierarchy level

Issue with Concurrent Execution in Regions



What is openValve set to?

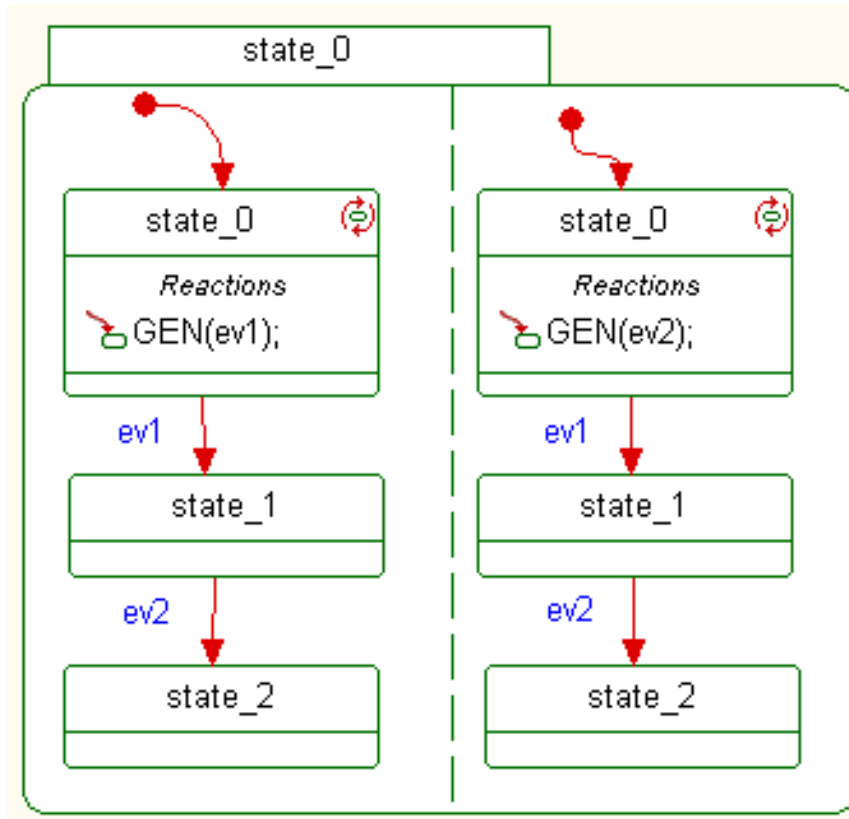
In ModelicaML: To "false" because *Region_0* was executed first and then *Region_1* is execute.

Priority Order

StateMachine_0::Region_0
StateMachine_0::Region_1



Issues With Concurrency When Using Event Queues (1/4)



-State machine enters the state_0 in both regions

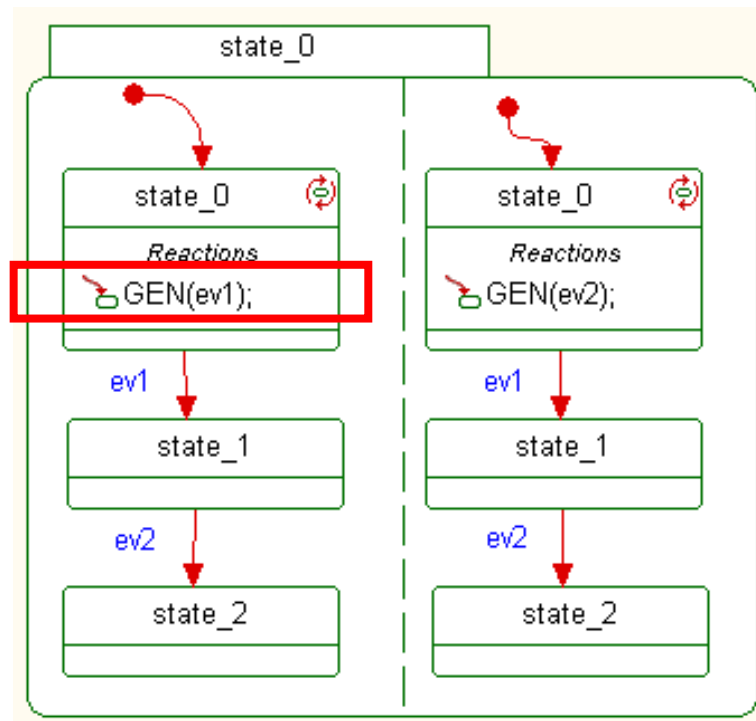
-On entry ev1 and ev2 are generated in state_0 in both regions

-Transitions to state_1 is performed in both regions

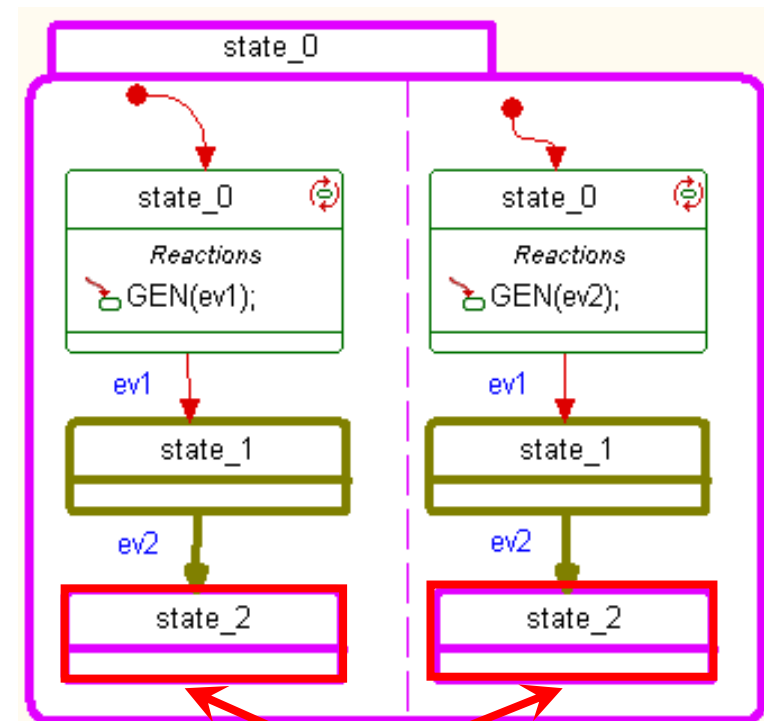
-...

What are the resulting active states configuration?

Issues With Concurrency When Using Event Queues (2/4)



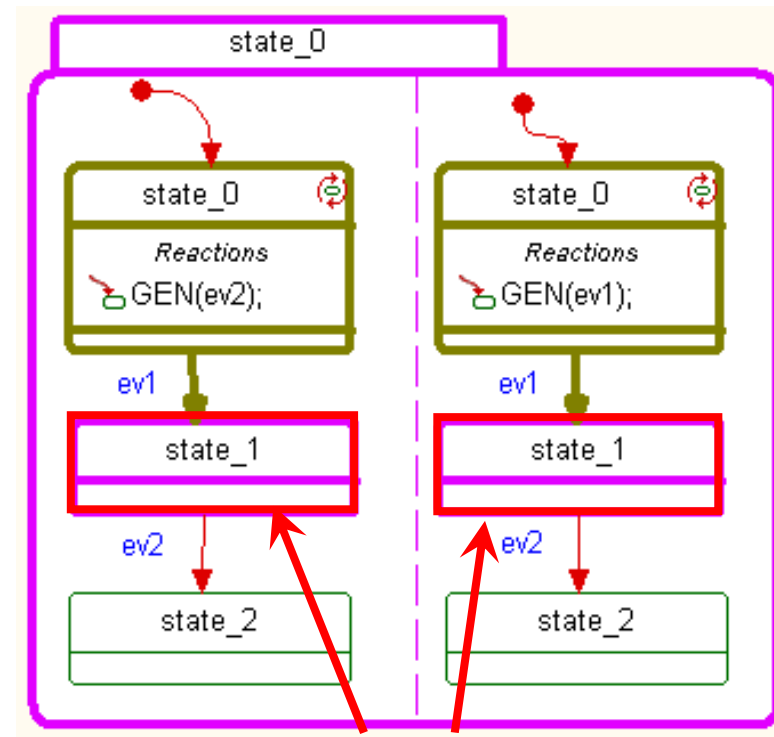
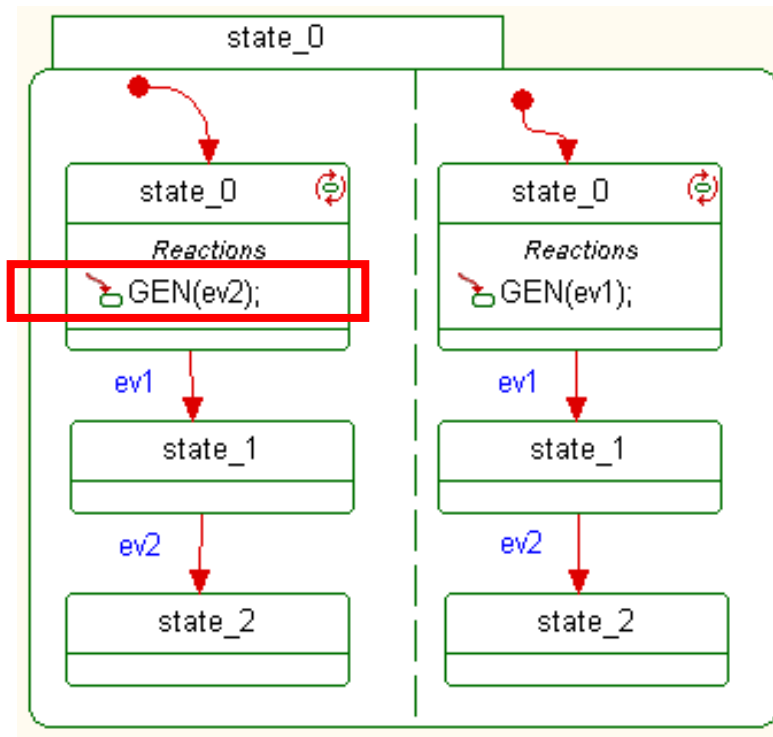
Simulation result in IBM Rhapsody



Active (resulting) states

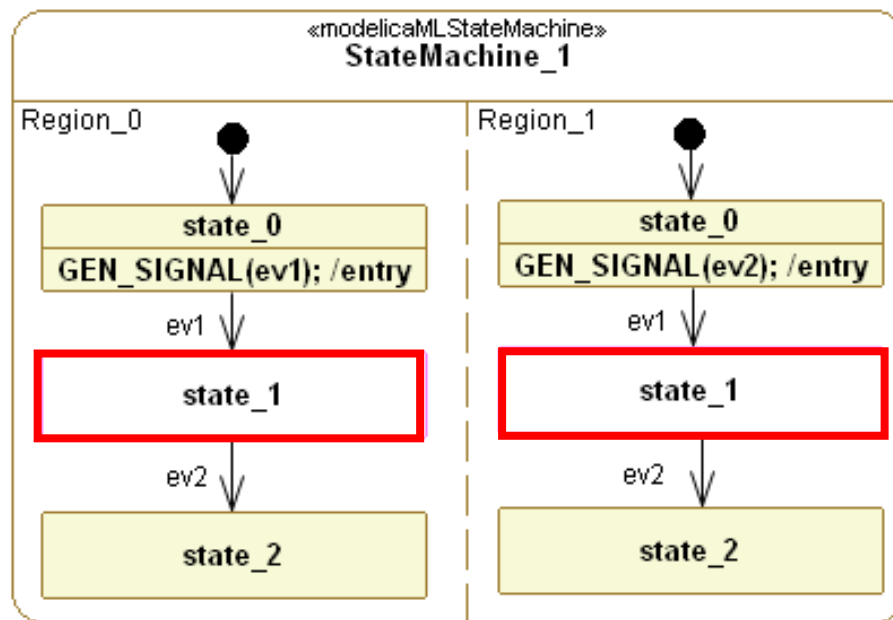
Issues With Concurrency When Using Event Queues (3/4)

Simulation result in IBM Rhapsody



Active (resulting) states

Issues With Concurrency When Using Event Queues (4/4)



Priority Order

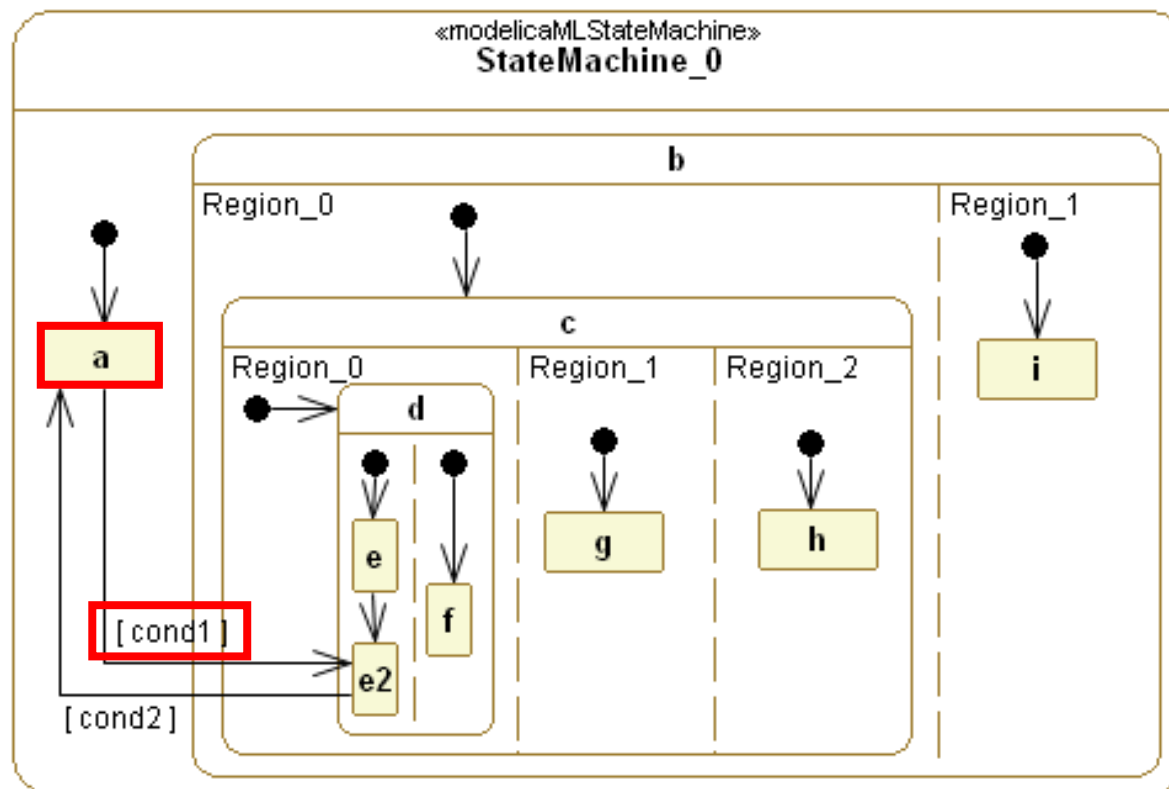
StateMachine_1::Region_0
 StateMachine_1::Region_1



What are the resulting states?

In ModelicaML: *Independent of whether ev1 is generated in Region_0 or in Region_1 the state machine ends up in state_1 in both regions because both events are generated when the state machines is in state_0 in both regions and are both consumed after the transition to state_1 is performed.*

Issues with Inter-Level Transitions

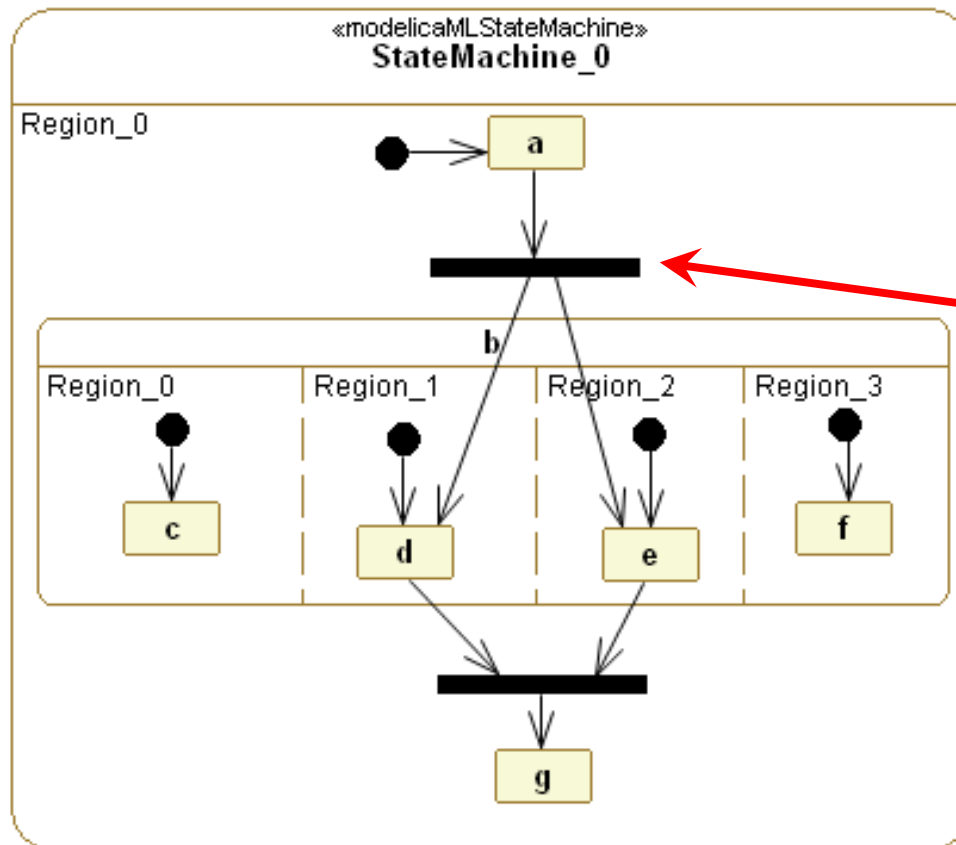


In which order are the states activated when **cond1** becomes true?

b, c, d, e2, i, h, g, f
b, c, d, e2, i, g, h, f
b, c, d, e2, f, g, h, i
b, c, d, e2, f, h, g, i

In ModelicaML:
b, c, d, e2, f, g, h, i

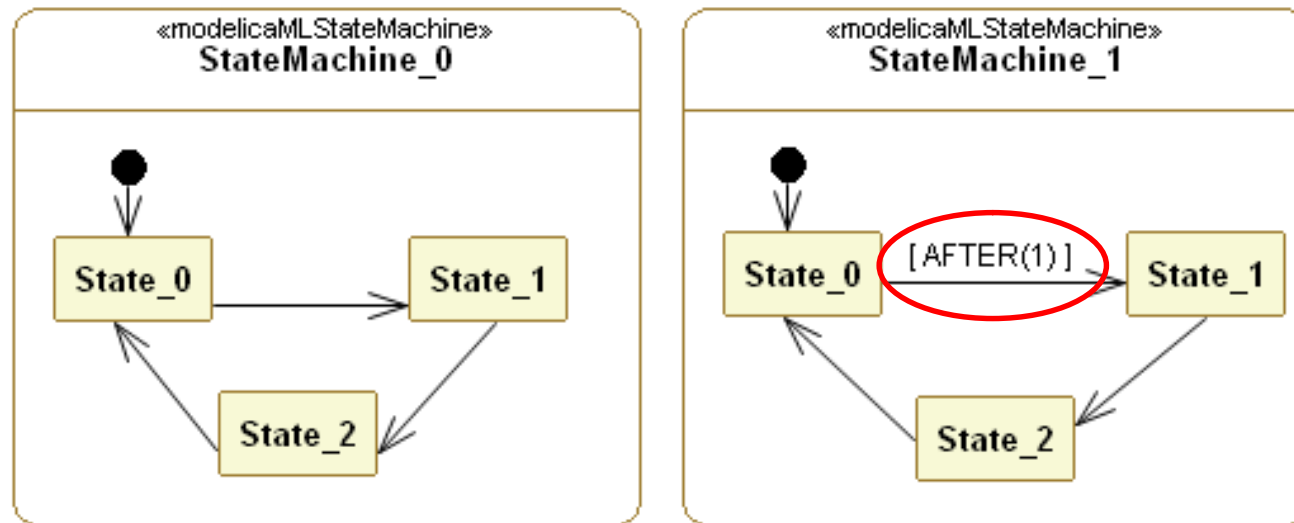
Issues with Fork and Join



In which sequence are states *b*, *c*, *d*, *e*, and *f* activated when the transitions (fork construct) from state *a* is executed?

In ModelicaML: *b*, *d* and *e* (based on the fork-outgoing transitions priority), *c* and *f* (based on their region priority)

Issues with Instantaneous States: Deadlocks (Infinite Looping)



Time-delayed transitions for breaking infinite looping at the same time instant.

Conclusion

- Using Modelica as execution language it is possible to support a comprehensive set of UML state machines specification
- Suggested improvement of the UML state machines specification
 - Priority for conflicting state-outgoing transitions
 - Priority for regions
- This enhancement of specification will clarify semantic and ensure
 - That state machine behavior is deterministic
 - That state machine behaves as intended by the modeler



Thank you for your attention!

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